

# Appendix B: Implementation Showcase

of the doctoral thesis: Unfolding semantics of the untyped  $\lambda$ -calculus with `letrec`

Jan Rochel <jan@rochel.info>

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To demonstrate the realisability of our method, and for further illustration, we include the output of our implementation for some examples from the thesis. The implementation is called `maxsharing` and is available on Hackage. It is written in Haskell and therefore requires the Haskell Platform to be installed. Then, `maxsharing` can be installed via `cabal-install` using the commands `cabal update` and `cabal install maxsharing` from the terminal. Invoke the executable `maxsharing` in your `cabal-directory` with a file as an argument that contains a  $\lambda_{\text{letrec}}$ -term. Run `maxsharing -h` for help on run-time flags.

Consider the following explanations for understanding the tool's output.

*$\lambda$ -letrec-term*: The original term as recognised by the parser

*translation used*: The user can specify which term graph semantics shall be used in processing the term. All further output is with respect to that translation. The two options are:

- $\llbracket \cdot \rrbracket_{\mathcal{T}}$  is called: `maximal prefix lengths but eager scope`
- $\llbracket \cdot \rrbracket_{\mathcal{T}}^{\text{min}}$  is called: `minimal prefix lengths`.

*scoped (adbmals)*: The term enriched by `adbmals`. The `adbmals` ( $\wedge$ ) is to be read as a scope delimiter that explicitly includes the name of the  $\lambda$ -variable whose scope it delimits. The `adbmals` are placed in accordance to the translation used.

*scoped (nameless)*: A nameless scoped representation, where the names of variables are omitted for abstractions as well as for abstraction variable occurrences, shown as a `0`-symbol. The scoping is expressed by scope-delimiters in the shape of an `S`-symbol. Note that these terms can be obtained by a simple syntactical transformation of the `adbmals`-terms.

*derivation*: derivation according to the proof system in fig. 3.4 that shows the translation as a stepwise process and includes all subterms with their abstraction prefixes. Note that in the notation of the prefixed terms the binding annotation of a variable is omitted if it is empty. The leftmost variable  $\star$  is omitted entirely if it has no binding annotations. Even though the correspondence between the derivations in the proof system in fig. 3.4 and the translations  $\llbracket \cdot \rrbracket_{\mathcal{T}}$  and  $\llbracket \cdot \rrbracket_{\mathcal{T}}^{\text{min}}$  is not provided here, we think that the derivation can help as an illustration of the translation process and its result.

*DFA*: The implementation produces a graphical depiction for the term's  $\lambda$ -term-graph in DFA form, as well as for the minimised form of the DFA; the pictures are included next to the textual output.

*spanning tree*: A textual representation of the minimised DFA's spanning tree, used for readback. It is displayed in first-order term rewriting notation, i.e. with a unary function symbols `L`, `S` for abstraction and scope delimiters, a binary function symbol `A` for application and a nullary symbol `0` for abstraction variable occurrences. Furthermore there is a class of function symbols written as a vertical bar followed by an upper-case variable name `|F`, `|G`, etc. which signify a vertex with multiple incoming non-backlink edges, and therefore a vertex that will be the root of a shared subterm. There is also a class of corresponding nullary function symbols `F`, `G`, etc. which represent non-backlink, non-spanning-tree edges to these shared vertices.

*readback*: readback of the minimised DFA.

§ B.1 (example 3.3.2). For this example only  $[\cdot]_{\mathcal{T}}^{\min}$  is shown since both semantics yield the exact same output.

---

$\lambda$ -letrec-term:

$(\lambda x. x) (\lambda x. x)$

translation used: minimal prefix lengths

scoped (adbmals):

$(\lambda x. x) (\lambda x. x)$

scoped (nameless):

$(\lambda. 0) (\lambda. 0)$

derivation:

----- 0	----- 0	
(x) x	(x) x	
----- $\lambda$	----- $\lambda$	
() $\lambda x. x$	( ) $\lambda x. x$	
----- @		
() $(\lambda x. x)$	$(\lambda x. x)$	

DFA: writing to file

spanning tree:

@(L(0), L(0))

readback:

$(\lambda x. x) (\lambda y. y)$

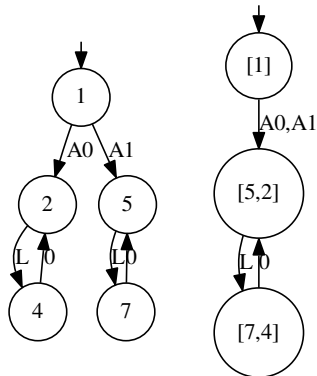
minimised DFA: writing to file

spanning tree:

@(|F(L(0)), F)

readback:

let F =  $\lambda x. x$   
in F F



$\lambda$ -letrec-term:

```

λf. let r = f r
    in r
    
```

translation used: minimal prefix lengths

scoped (adbmals):

```

λf. let r = f r
    in r
    
```

scoped (nameless):

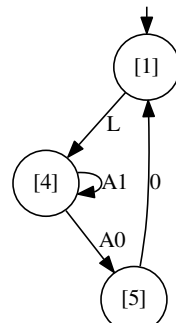
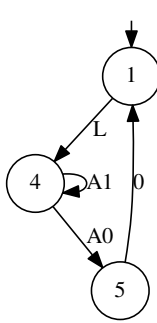
```

λ. let r = 0 r
    in r
    
```

derivation:

```

----- 0
(f{r}) f   (f{r}) r
-----  @
(f{r}) f r   (f{r}) r
-----  let
(f) let r = f r
    in r
-----  λ
() λf. let r = f r
    in r
    
```



DFA: writing to file

spanning tree:

L(|F(@0, F))

readback:

```

λx. let F = x F
    in F
    
```

minimised DFA: writing to file

spanning tree:

L(|F(@0, F))

readback:

```

λx. let F = x F
    in F
    
```

$\lambda$ -letrec-term:

```

λf. let r = f (f r)
    in r
    
```

translation used: minimal prefix lengths

scoped (adbmals):

```

λf. let r = f (f r)
    in r
    
```

scoped (nameless):

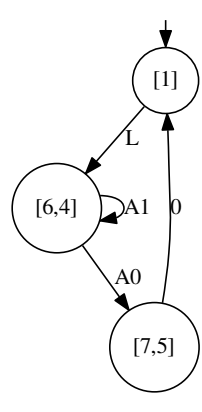
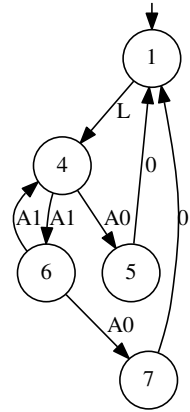
```

λ. let r = 0 (0 r)
    in r
    
```

derivation:

```

----- 0
(f{r}) f   (f{r}) r
-----  @
(f{r}) f   (f{r}) f r
-----  @
(f{r}) f (f r)   (f{r}) r
-----  let
(f) let r = f (f r)
    in r
-----  λ
() λf. let r = f (f r)
    in r
    
```



DFA: writing to file

spanning tree:

L(|F(@0, @0, F)))

readback:

```

λx. let F = x (x F)
    in F
    
```

minimised DFA: writing to file

spanning tree:

L(|F(@0, F))

readback:

```

λx. let F = x F
    in F
    
```

$\lambda$ -letrec-term:  
 let f =  $\lambda x. (\lambda y. f y) x$   
 in f

translation used: minimal prefix lengths

scoped (adbmals):  
 let f =  $\lambda x. /x. \lambda y. /y. f y x$   
 in f

scoped (nameless):  
 let f =  $\lambda. S((\lambda. S(f) 0)) 0$   
 in f

derivation:  
 $(\{f\}) f$   
 $\frac{}{(\{f\}) f} S \frac{}{0} 0$   
 $(\{f\} y) f \quad (\{f\} y) y$   
 $\frac{}{(\{f\} y) f y} @$   
 $(\{f\}) \lambda y. f y$   
 $\frac{}{(\{f\}) \lambda y. f y} S \frac{}{0} 0$   
 $(\{f\} x) \lambda y. f y \quad (\{f\} x) x$   
 $\frac{}{(\{f\} x) (\lambda y. f y) x} @$   
 $(\{f\}) \lambda x. (\lambda y. f y) x$   
 $\frac{}{(\{f\}) \lambda x. (\lambda y. f y) x} \lambda \frac{}{(\{f\}) f} let$   
 $(\{f\}) let f = \lambda x. (\lambda y. f y) x$   
 $\quad in f$

DFA: writing to file

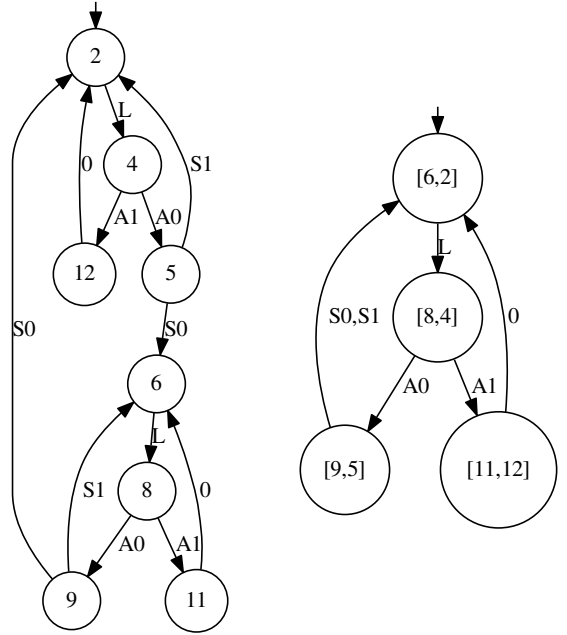
spanning tree:  
 $|F(L(@S(L(@S(F), 0))), 0))$

readback:  
 let F =  $\lambda x. /x. \lambda y. /y. F y x$   
 in F

minimised DFA: writing to file

spanning tree:  
 $|F(L(@S(F), 0))$

readback:  
 let F =  $\lambda x. F x$   
 in F



$\lambda$ -letrec-term:  
 let f =  $\lambda x. (\lambda y. f x) x$   
 in f

translation used: minimal prefix lengths

scoped (adbmals):  
 let f =  $\lambda x. (\lambda y. /y. /x. f x) x$   
 in f

scoped (nameless):  
 let f =  $\lambda. (\lambda. S((S(f) 0))) 0$   
 in f

derivation:  
 $(\{f\}) f$   
 $\frac{}{(\{f\}) f} S \frac{}{0} 0$   
 $(\{f\} x) f \quad (\{f\} x) x$   
 $\frac{}{(\{f\} x) f x} @$   
 $(\{f\}) \lambda x. f x$   
 $\frac{}{(\{f\}) \lambda x. f x} S$   
 $(\{f\} x y) f x$   
 $\frac{}{(\{f\} x y) f x} \lambda \frac{}{0} 0$   
 $(\{f\} x) \lambda y. f x \quad (\{f\} x) x$   
 $\frac{}{(\{f\} x) (\lambda y. f x) x} @$   
 $(\{f\}) \lambda x. (\lambda y. f x) x$   
 $\frac{}{(\{f\}) \lambda x. (\lambda y. f x) x} \lambda \frac{}{(\{f\}) f} let$   
 $(\{f\}) let f = \lambda x. (\lambda y. f x) x$   
 $\quad in f$

DFA: writing to file

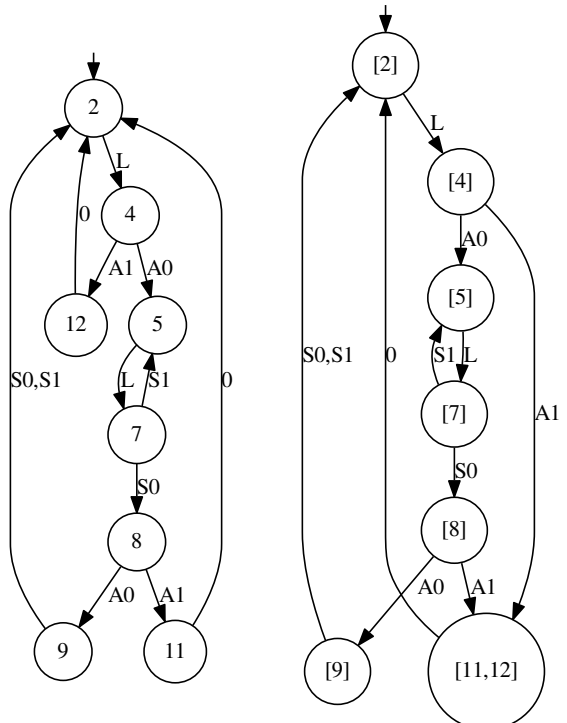
spanning tree:  
 $|F(L(@L(S(@S(F), 0))), 0))$

readback:  
 let F =  $\lambda x. (\lambda y. /y. /x. F x) x$   
 in F

minimised DFA: writing to file

spanning tree:  
 $|F(L(@L(S(@S(F), |G(0))), G))$

readback:  
 let F =  $\lambda x. let G = x$   
 $\quad in (\lambda y. F G) G$   
 in F



§ B.4 (fig. 3.6). Also for the term  $\lambda a. \lambda b. \text{let } f = a \text{ in } a a (f a) b$  from fig. 3.6 the translations are identical for  $\llbracket \cdot \rrbracket_{\mathcal{T}}$  and  $\llbracket \cdot \rrbracket_{\mathcal{T}}^{\min}$ .

$\lambda$ -letrec-term:  
 $\lambda a. \lambda b. \text{let } f = a$   
in  $a a (f a) b$

translation used: minimal prefix lengths

scoped (adbmals):  
 $\lambda a. \lambda b. \text{let } f = a$   
in  $/b. a a (f a) b$

scoped (nameless):  
 $\lambda. \lambda. \text{let } f = 0$   
in  $S((0 0 (f 0))) 0$

derivation:

```

----- 0 ----- 0 ----- 0
(a{f}) a (a{f}) a (a{f}) f (a{f}) a
----- @ ----- @
(a{f}) a a (a{f}) f a
----- @
(a{f}) a a (f a)
----- S ----- 0
(a{f} b) a a (f a) (a{f} b) b
----- @
(a{f}) a (a{f} b) a a (f a) b
----- let
(a b) let f = a
in a a (f a) b
----- λ
(a) λb. let f = a
in a a (f a) b
----- λ
() λa. λb. let f = a
in a a (f a) b

```

DFA: writing to file

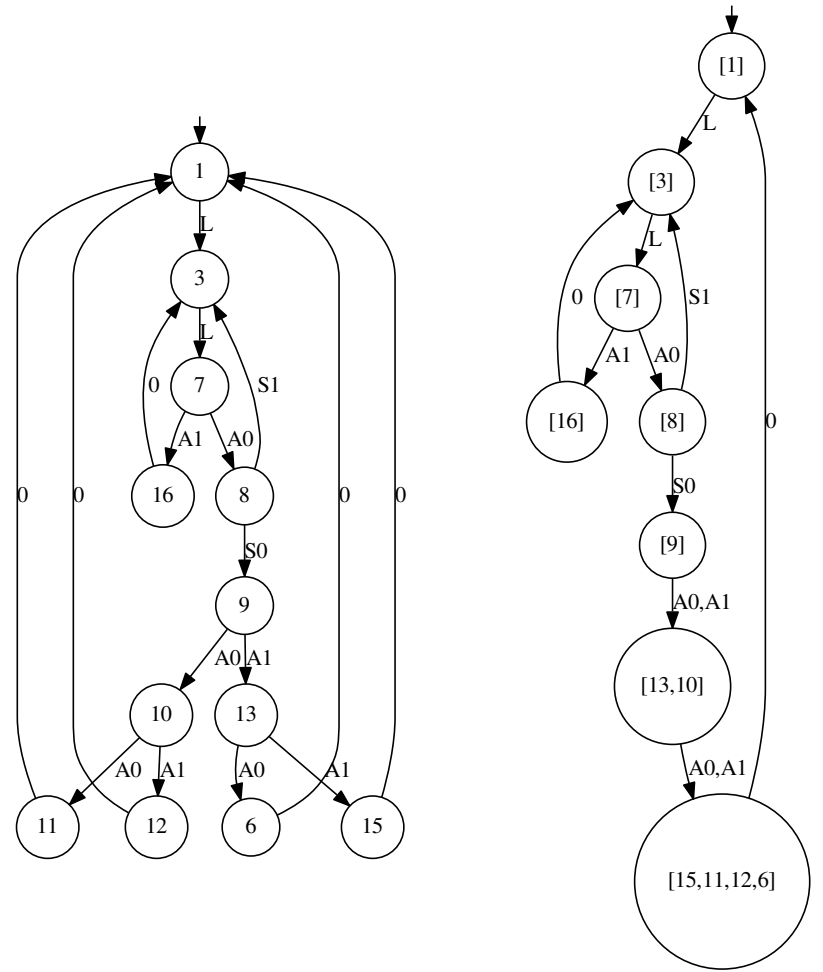
spanning tree:  
 $L(L(@(@(@(@(@(@(0, 0)), @(@(0, 0))), 0))), 0)))$

readback:  
 $\lambda x. \lambda y. /y. x x (x x) y$

minimised DFA: writing to file

spanning tree:  
 $L(L(@(@(@(|F@(|G(0), G)), F)), 0)))$

readback:  
 $\lambda x. \text{let } F = G G$   
 $G = x$   
in  $\lambda y. F F y$



To illustrate the necessity for different prefix lengths (§ 3.5.17) we also provide the translation with equal abstraction-prefix lengths, which is not eager scope.

$\lambda$ -letrec-term:  
 $\lambda a. \lambda b. \text{let } f = a$   
in  $a a (f a) b$

translation used: maximal prefix lengths

scoped (adbmals):  
 $\lambda a. \lambda b. \text{let } f = /b. a$   
in  $/b. a a (f /b. a) b$

scoped (nameless):  
 $\lambda. \lambda. \text{let } f = S(0)$   
in  $S((0 0)) (f S(0)) 0$

```

----- 0 ----- 0 ----- 0
(a) a (a) a (a) a
----- @ ----- S
(a) a a (a b{f}) f (a b{f}) a
----- S ----- @
(a b{f}) a a (a b{f}) f a
----- @ ----- 0
(a) a (a b{f}) a a (f a) (a b{f}) b
----- S ----- @
(a b{f}) a (a b{f}) a a (f a) b
----- let
(a b) let f = a
in a a (f a) b
----- λ
(a) λb. let f = a
in a a (f a) b
----- λ
() λa. λb. let f = a
in a a (f a) b

```

DFA: writing to file

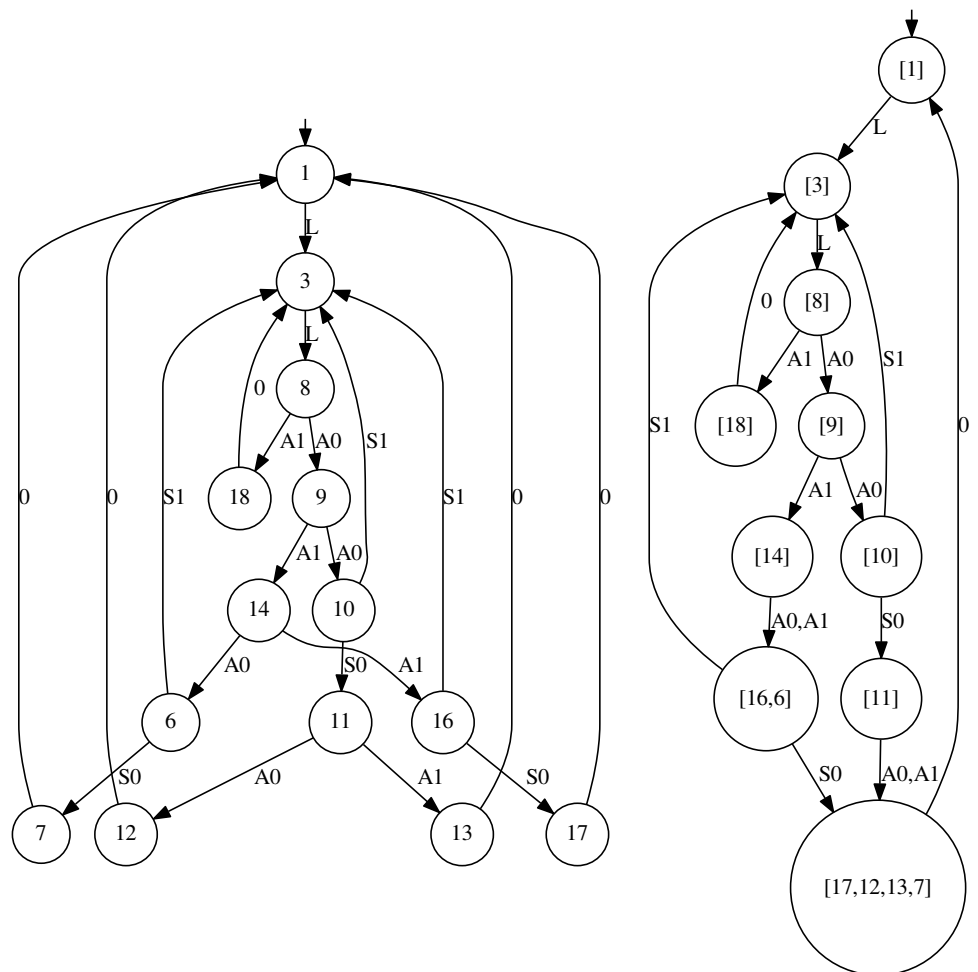
spanning tree:  
 $L(L(@(@(@(@(@(@(0, 0))), @(@S(0), S(0))), 0))), 0)))$

readback:  
 $\lambda x. \lambda y. /y. x x (/y. x /y. x) y$

minimised DFA: writing to file

spanning tree:  
 $L(L(@(@(@(|G(0), G)), @(|F(S(G)), F)), 0)))$

readback:  
 $\lambda x. \text{let } G = x$   
in  $\lambda y. \text{let } F = G$   
in  $G G (F F) y$





```

λ-letrec-term:
λx. λy. let I = λz. z
      f = x
      in y I (I y) (f f)
  
```

translation used: maximal prefix lengths but eager scope

```

scoped (adbmals):
λx. λy. let I = /y. /x. λz. z
      f = x
      in y I (I y) /y. f f
  
```

```

scoped (nameless):
λ. λ. let I = S(S(λ. 0))
      f = 0
      in 0 I (I 0) S((f f))
  
```

```

derivation:
----- 0
(z) z
----- λ
() λz. z
----- S
(x{f}) λz. z
----- S
(x{f} y{I}) λz. z
----- S
(x y) let I = λz. z
      f = x
      in y I (I y) (f f)
----- λ
(x) λy. let I = λz. z
      f = x
      in y I (I y) (f f)
----- λ
() λx. λy. let I = λz. z
      f = x
      in y I (I y) (f f)
  
```

DFA: writing to file

```

spanning tree:
L(L(@(@(@0, |F(S(S(L(0))))), @F, 0)), S(@(|G(0), G))))
  
```

```

readback:
λx. let G = x
      in λy. let F = /y. /x. λz. z
            in y F (F y) /y. G G
  
```

minimised DFA: writing to file

```

spanning tree:
L(L(@(@(|H(0), |F(S(S(L(0))))), @F, H)), S(@(|G(0), G))))
  
```

```

readback:
λx. let G = x
      in λy. let H = y
            F = λz. z
            in H F (F H) (G G)
  
```

```

λ-letrec-term:
λx. λy. let I = λz. z
      f = x
      in y I (I y) (f f)
  
```

translation used: minimal prefix lengths

```

scoped (adbmals):
λx. λy. let I = λz. z
      f = x
      in y /y. /x. I (/y. /x. I y) /y. f f
  
```

```

scoped (nameless):
λ. λ. let I = λ. 0
      f = 0
      in 0 S(S(I)) (S(S(I)) 0) S((f f))
  
```

```

derivation:
----- 0
(*{I} z) z
----- λ
(*{I}) λz. z
----- S
(x y) let I = λz. z
      f = x
      in y I (I y) (f f)
----- λ
(x) λy. let I = λz. z
      f = x
      in y I (I y) (f f)
----- λ
() λx. λy. let I = λz. z
      f = x
      in y I (I y) (f f)
  
```

DFA: writing to file

```

spanning tree:
L(L(@(@(|H(0), |G(S(S(L(0))))), @S(S(F)), 0)), S(@(|G(0), G))))
  
```

```

readback:
let F = λz. z
in λx. let G = x
      in λy. y /y. /x. F (/y. /x. F y) /y. G G
  
```

minimised DFA: writing to file

```

spanning tree:
L(L(@(@(|H(0), |G(S(S(L(0))))), @G, H)), S(@(|F(0), F))))
  
```

```

readback:
λx. let F = x
      in λy. let H = y
            G = λz. z
            in H G (G H) (F F)
  
```

